

# Arbitrageurs' profits, LVR, and sandwich attacks: batch trading as an AMM design response

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Blockchain@X - OMI Workshop

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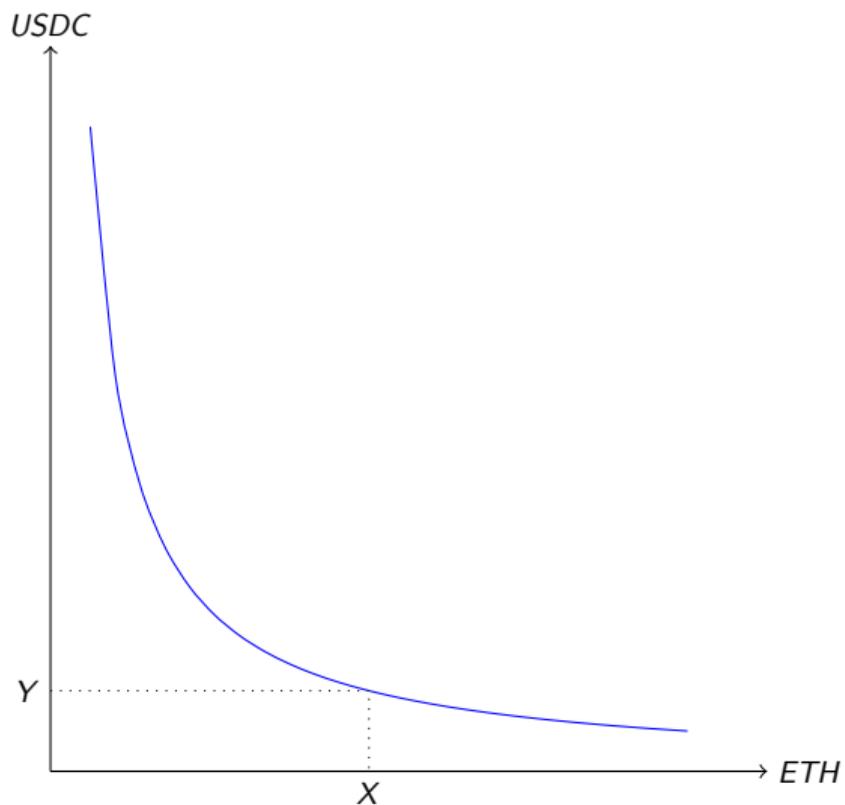
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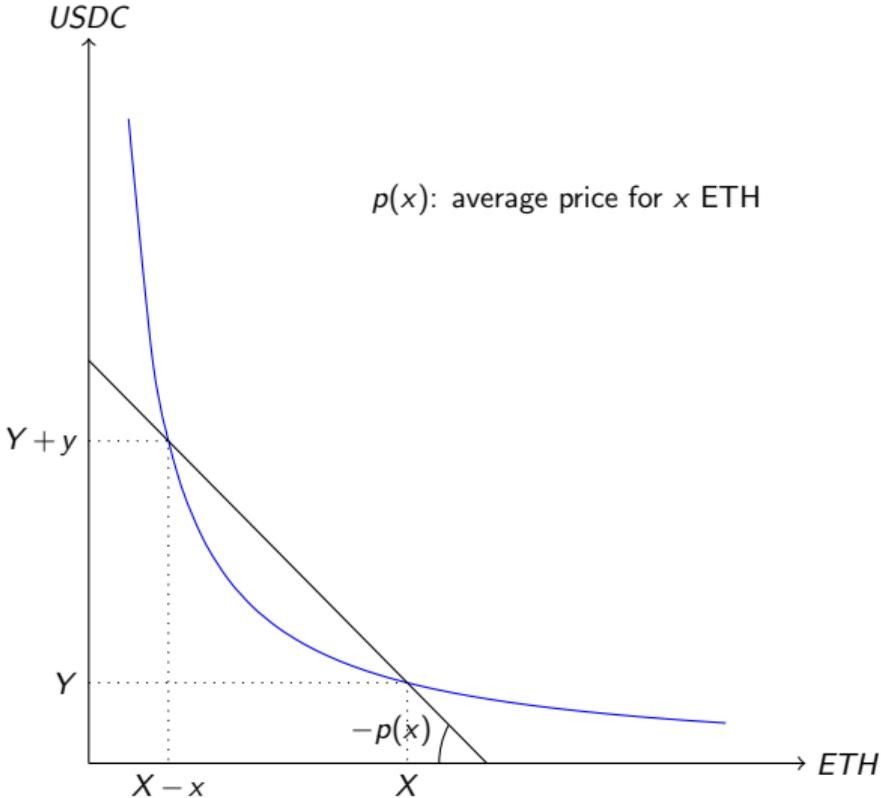
- The invariant function must satisfy path-independence
- Most common invariant function: the product  $\Rightarrow p(x) = y/x = Y/(X - x)$

# Constant Product Automated Market Maker

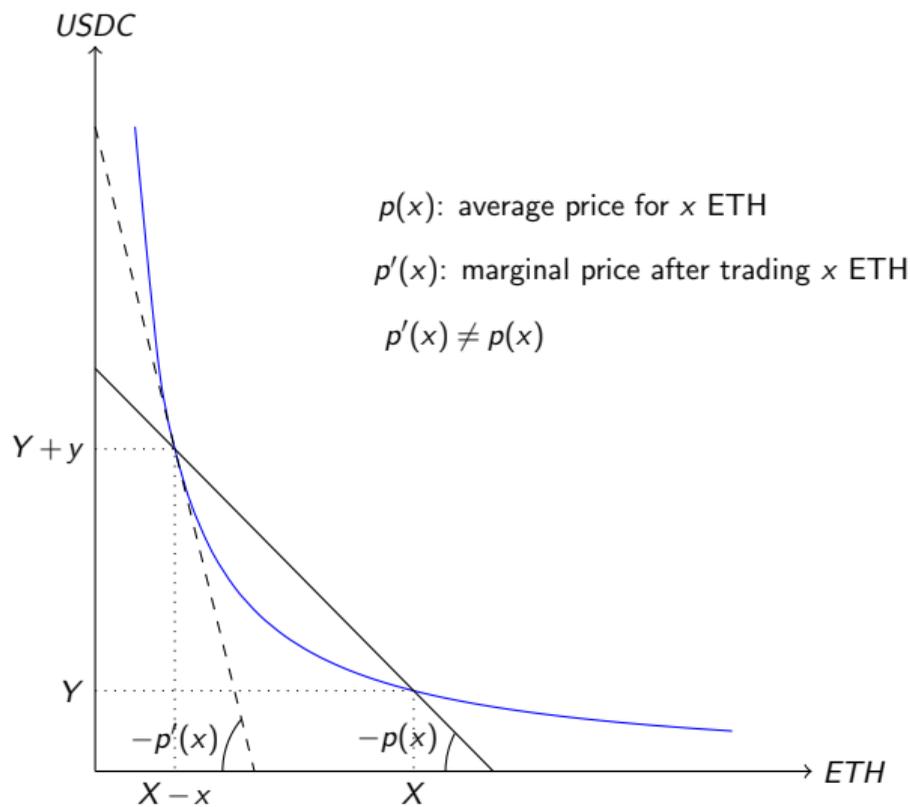




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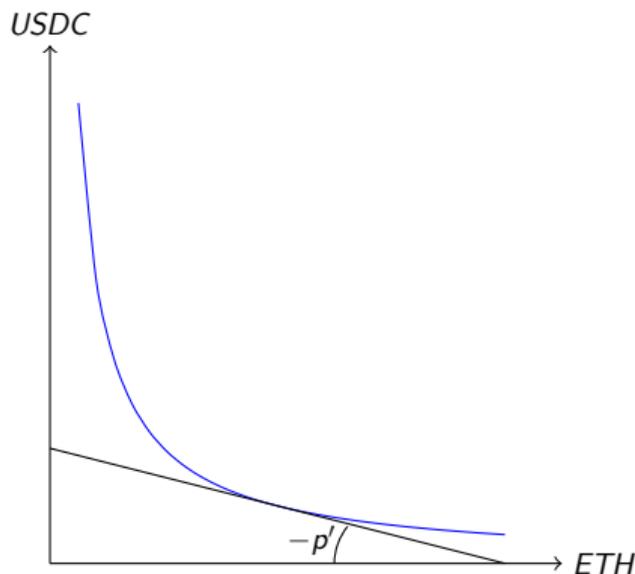


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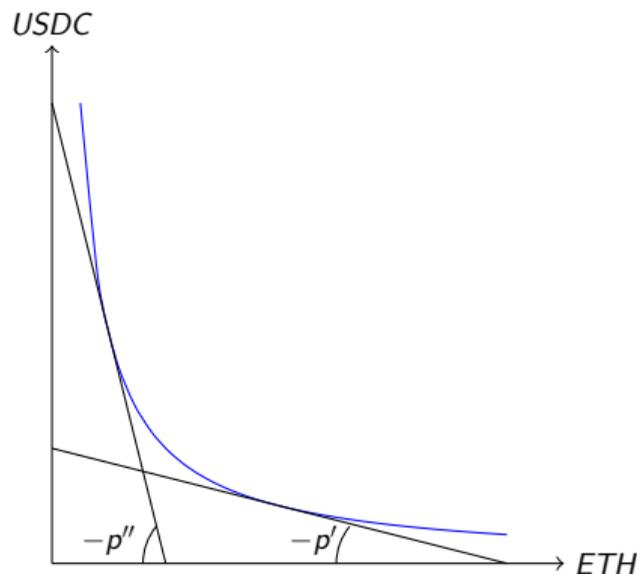
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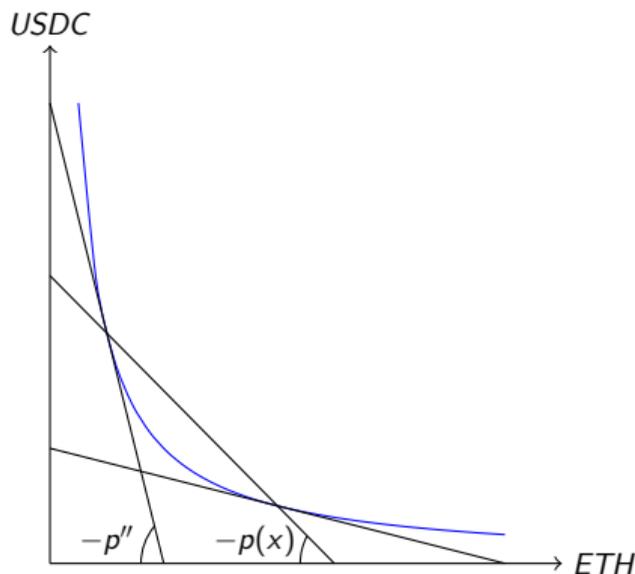
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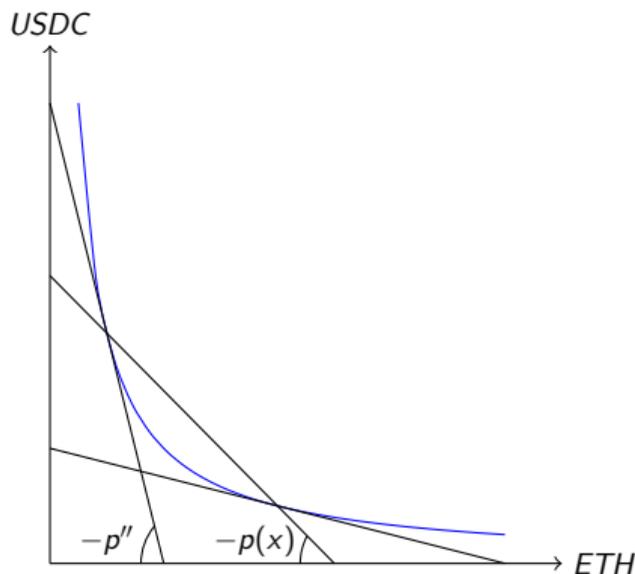
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- Defensive mechanism: fees



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- back run the victim with the opposite trade (sell ETH)
- $\Rightarrow$  the attacker buys cheap and sells expensive; the victim buys expensive

## Our paper: batching trades (and designing the AMM around this) solves both problems

- All trades are collected off-chain and batched:
  - ▶ they are settled p2p if possible; the remaining is settled on an AMM (at the same price)
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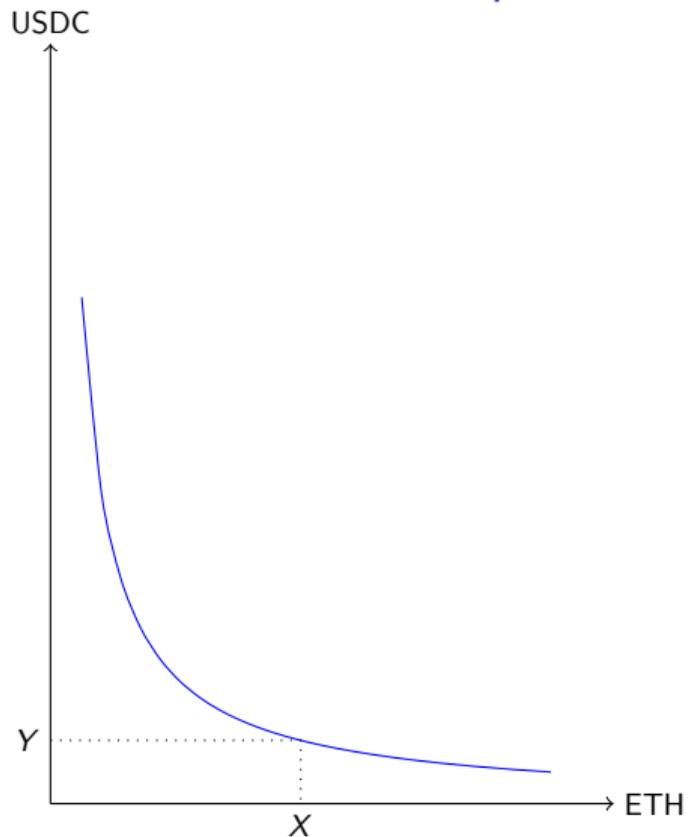
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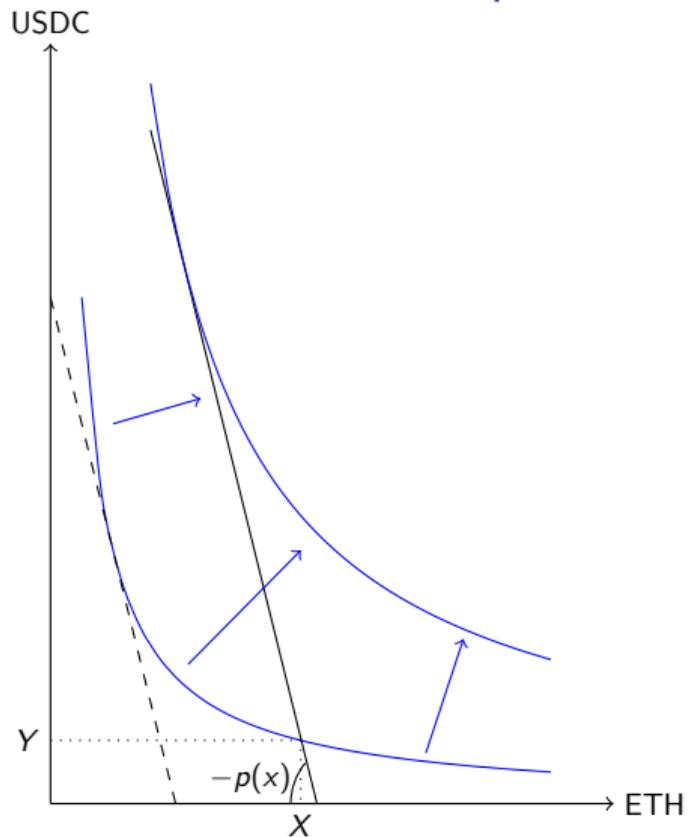
Hence, to purchase  $x$  ETH on the FM-AMM, the price needs to be:

$$p^{FM-AMM}(x) = \frac{Y}{X - 2x}.$$

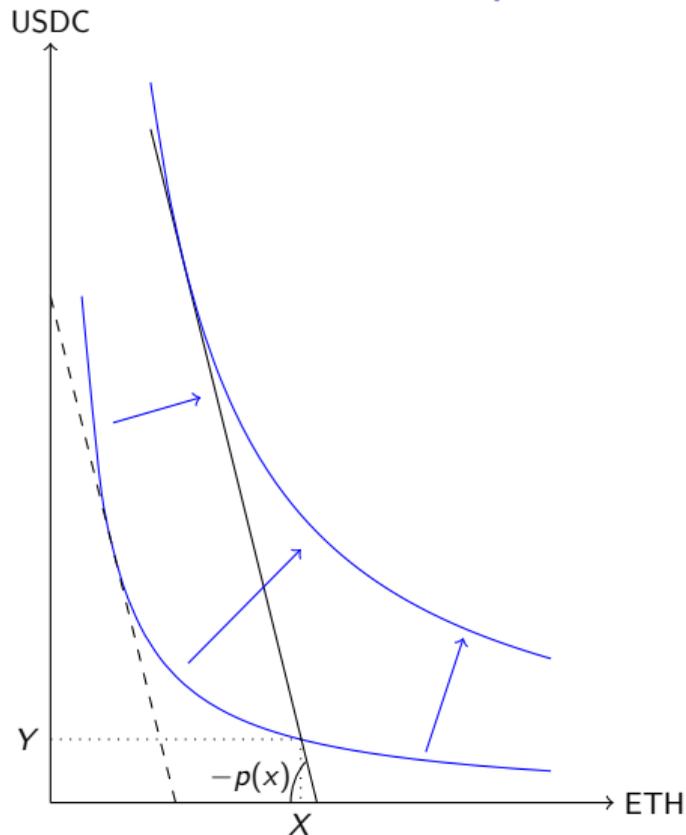
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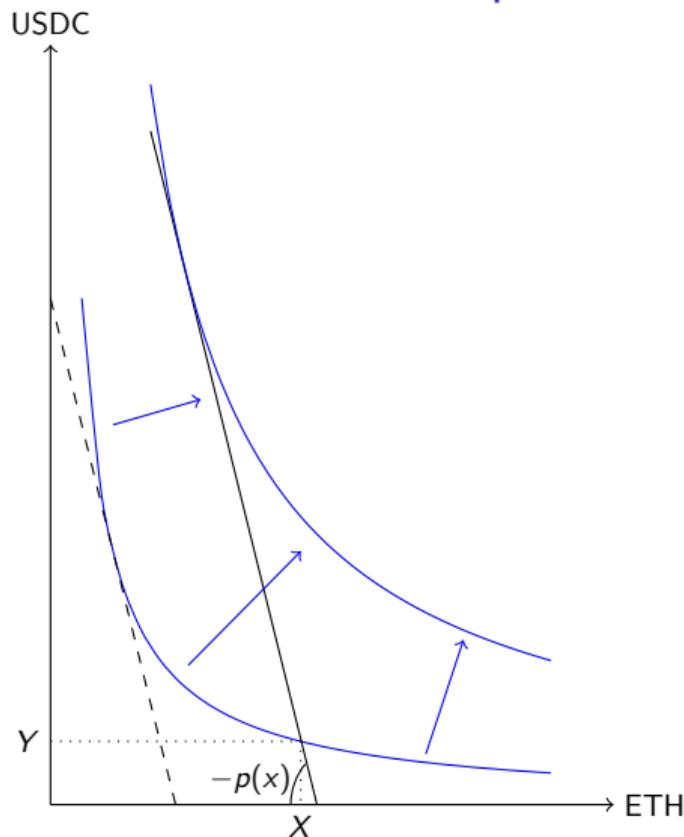


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- FM-AMM is **clearing-price consistent**: the price at which it trades equals the marginal price after the trade
- FM-AMM violates path independence: it can be exploited by splitting trades → **batching is necessary**

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  - ▶ *Even periods* = different block (on-chain transaction occurs)
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- FM-AMM charges no fee for inclusion in a batch, and a fee  $\tau$  (in the input token) for settling an order on the FM-AMM

## Theoretical model: FM-AMM in equilibrium

### Proposition:

Suppose that, at the end of an even period, the reserves of the FM-AMM are  $X$  and  $Y$ . In the equilibrium of the subsequent odd period, after  $p^*$  is realized, if noise traders collectively submit trade  $A$  to the batch, then arbitrageurs will collectively submit trade  $y(p^*)$  such that

$$\tilde{p}(A + y(p^*), \tau) = p^*$$

where  $\tilde{p}(A + y(p^*), \tau)$  is the FM-AMM effective price (i.e., price after fees).

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- $p^*$  determines the rebalancing trade, which determines the fees earned

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AMM function		
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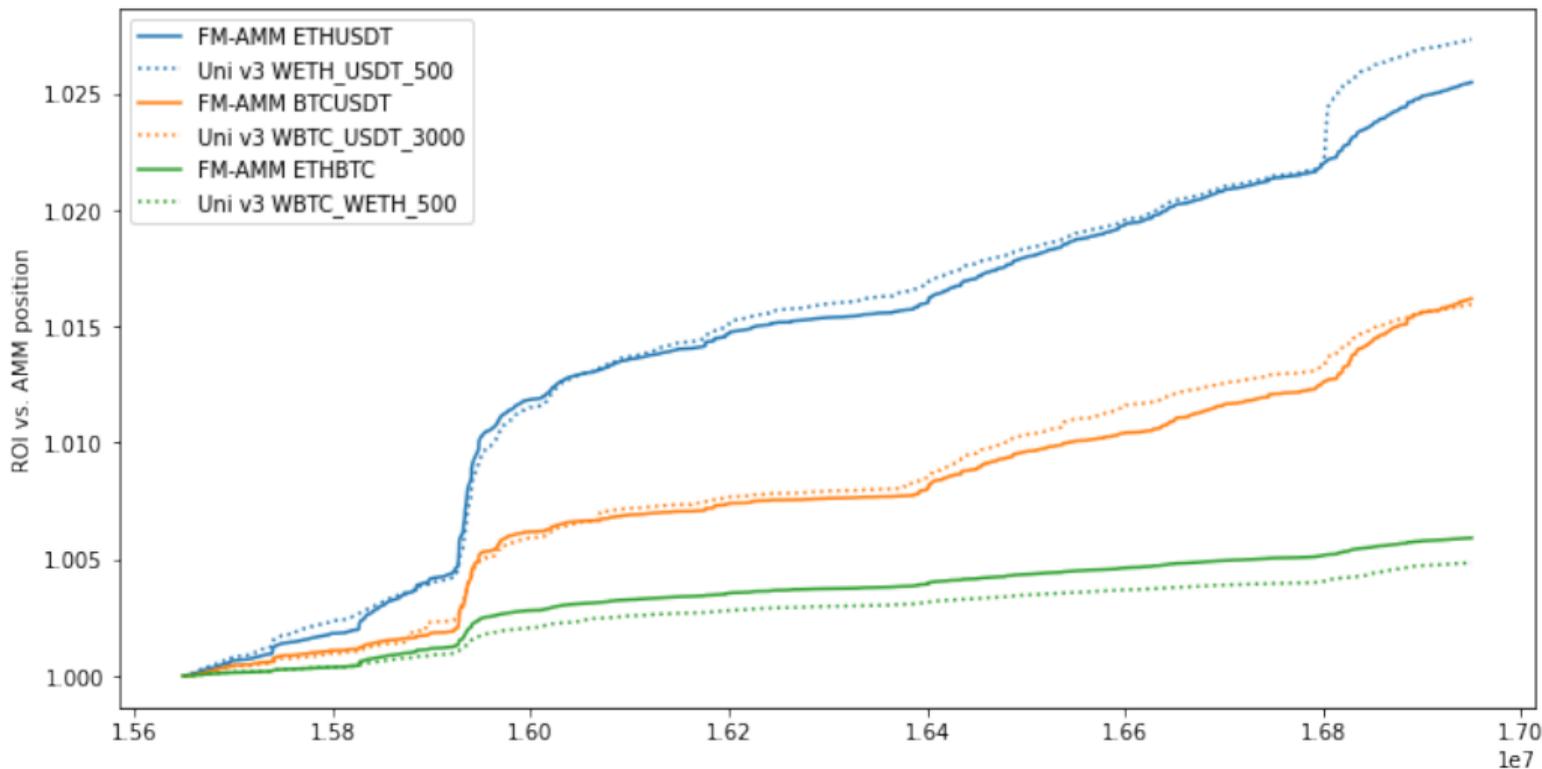
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- Compute the return of liquidity providers on Uniswap v3 (the leading AMM) over the same period for the pools WETH-USDT (with fee 0.05%), WBTC-USDT (with fee 0.3%), and WBTC-WETH (with fee 0.05%).
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- Compare the two.

# FM-AMM vs Uniswap v3 pool





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- **The differences in returns are small.**

# Conclusions

- Batching allows for a novel AMM design that eliminates arbitrageurs' profits (LVR) and sandwich attacks.
- (for the period and the token pair we consider) for liquidity providers, an FM-AMM that does not earn fees from noise traders performs as well as Uniswap v3
  - ▶ An FM-AMM that also earns fees from noise traders should perform better

# Literature

- Batch auction: Budish, Cramton, and Shim (2015)
- Sandwich attacks: Park (2022), Torres et al. (2021), Qin et al. (2022).
- Arb profits and LVR: Aoyagi (2020), Capponi and Jia (2021), and Milionis et al. (2022), Milionis et al. (2023)
- surplus maximizing AMM / axiomatization of AMM: Goyal et al. (2022), Schlegel and Mamageishvili (2022)
- several blog posts

Thank you!